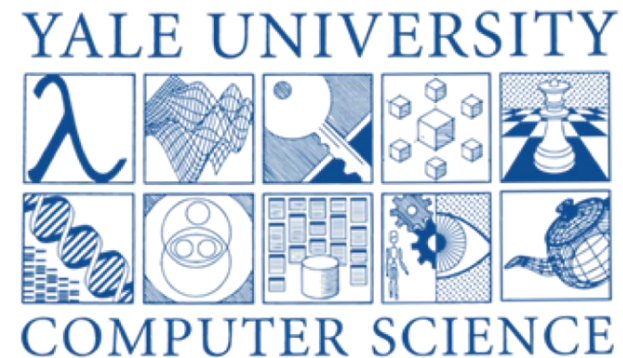


# Panic Recovery in Rust-based Embedded Systems

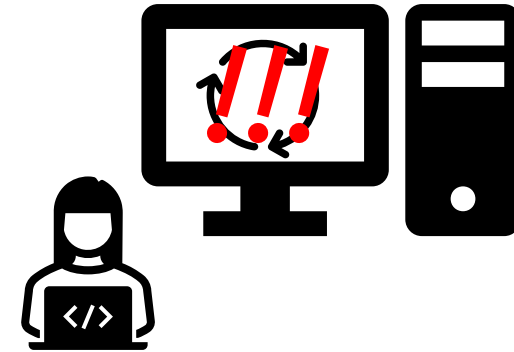
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Efficient Computing Lab  
Yale University



# Tasks May Crash Due to (Language) Exception

- Language exception
  - Rust panic, C++ `std::runtime_error`, etc.
- Various sources of exception
  - Bugs in program code
  - Failed assertions
  - Transient hardware error
- Embedded systems pose greater challenges
  - Unattended
  - Mission critical



# Recover by Unwinding & Restarting

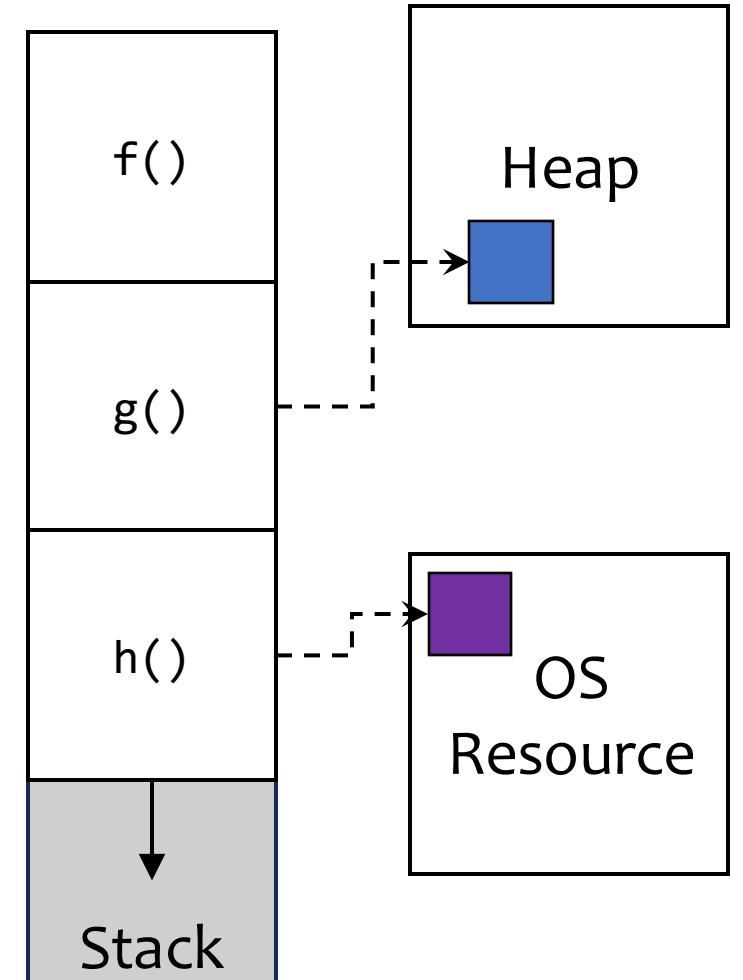
- Reclaim resources by unwinding a task's stack.
  - Force function returns out of `main()`.
  - Destruct initialized objects.
- Resume execution by restarting the task.
  - Run again from `main()`.
- Applicable to soft real-time systems.



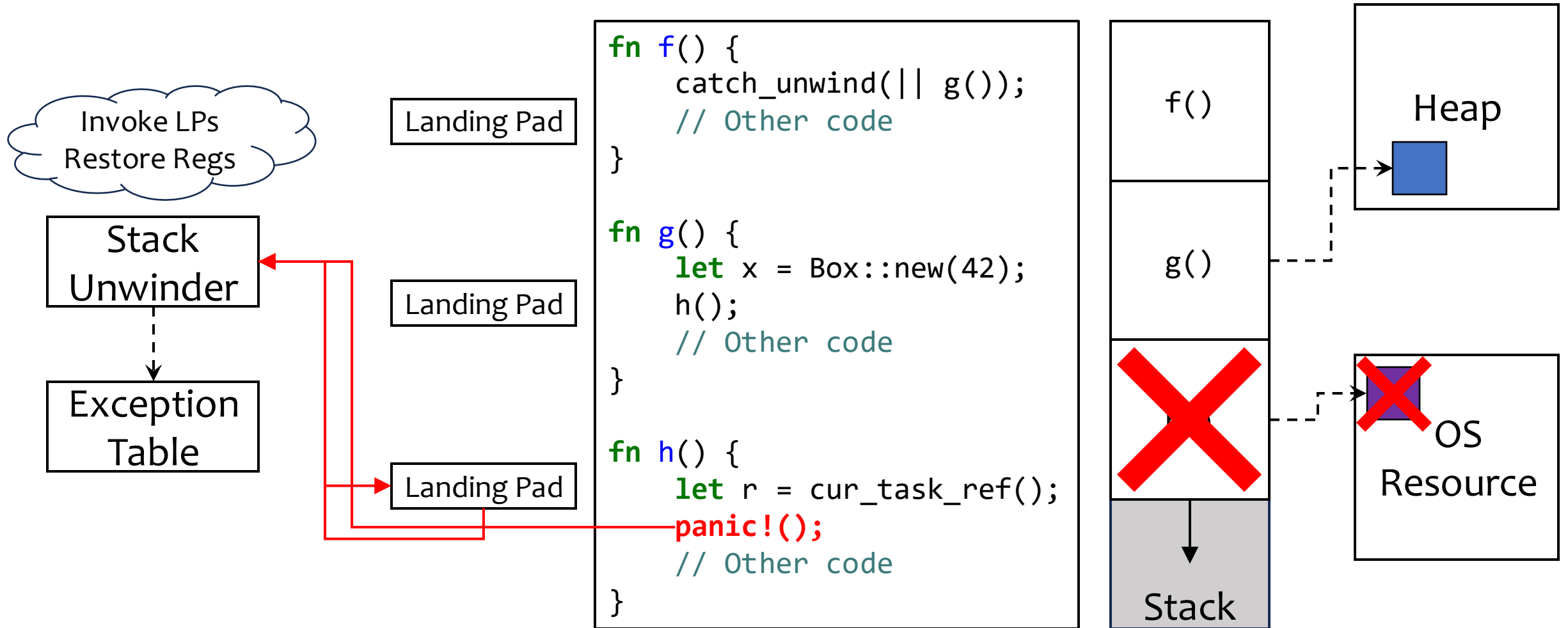
# How Stack Unwinding works

# Unwinder forces return until catch\_unwind

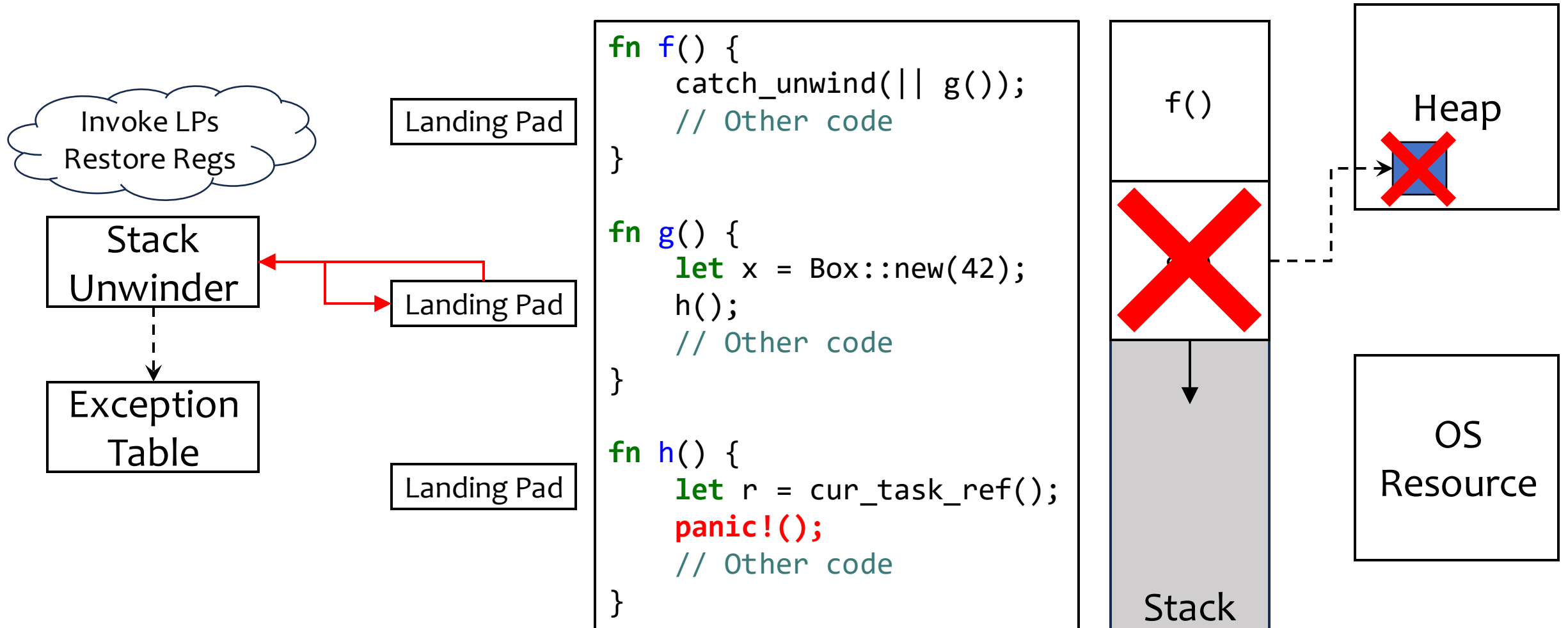
```
fn f() {  
    catch_unwind(|| g());  
    // Other code  
}  
  
fn g() {  
    let x = Box::new(42);  
    h();  
    // Other code  
}  
  
fn h() {  
    let r = cur_task_ref();  
    panic!();  
    // Other code  
}
```



# Unwinder forces return until catch\_unwind



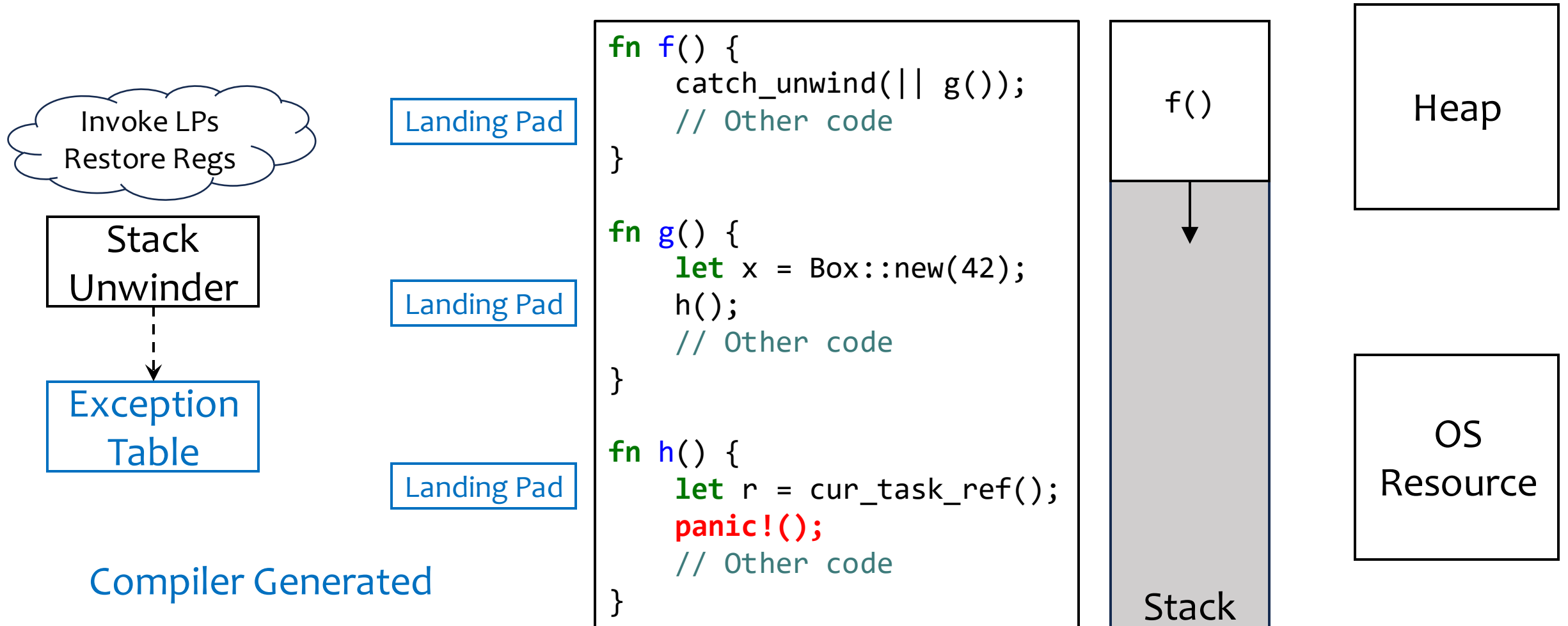
# Unwinder forces return until catch\_unwind







# Unwinder forces return until catch\_unwind



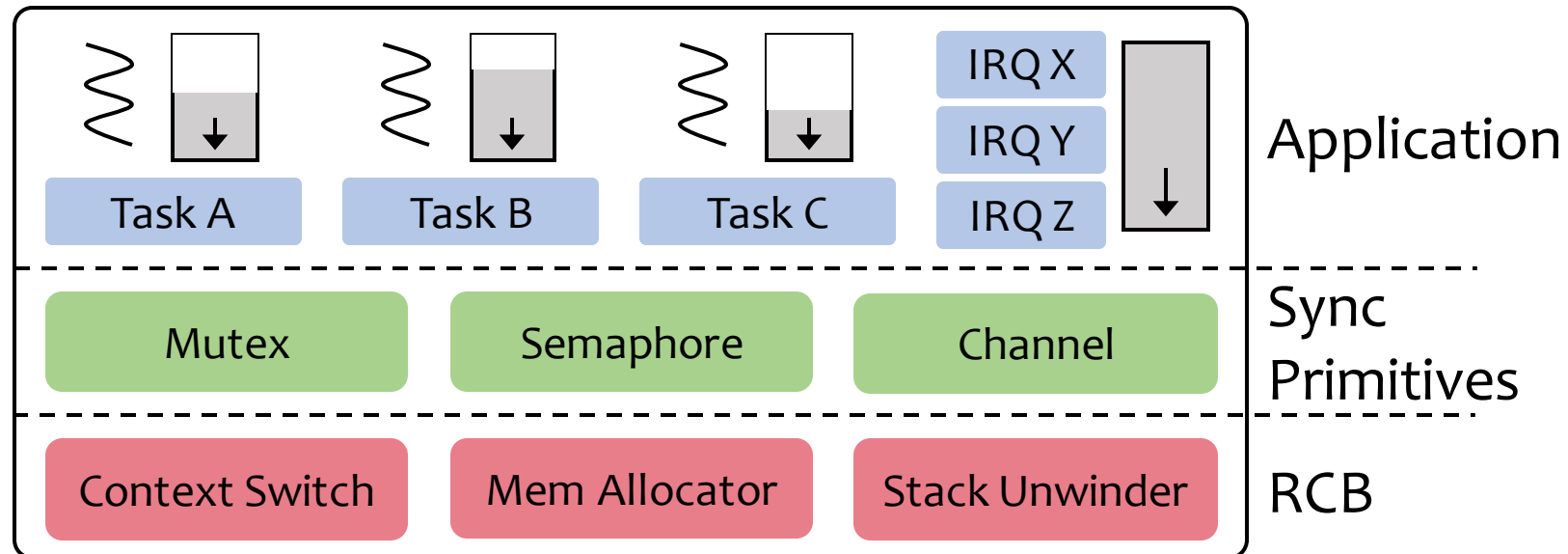
# Challenges of Unwinding on Embedded Systems

- Storage overhead
  - Unwinder logic, landing pads, exception table
  - Up to 5x size increase
  - Rust panic's simplicity: only `Exception` type, always fatal, no re-throw, etc.
  - Arm EHABI more compact exception table format
- Performance overhead
  - Unwinder interpreting exception table
  - Up to 1000x slow down
  - Concurrent restart-able task abstraction (facilitated by Rust)

# Hopter: An Embedded OS Capable of Recovering from Panic

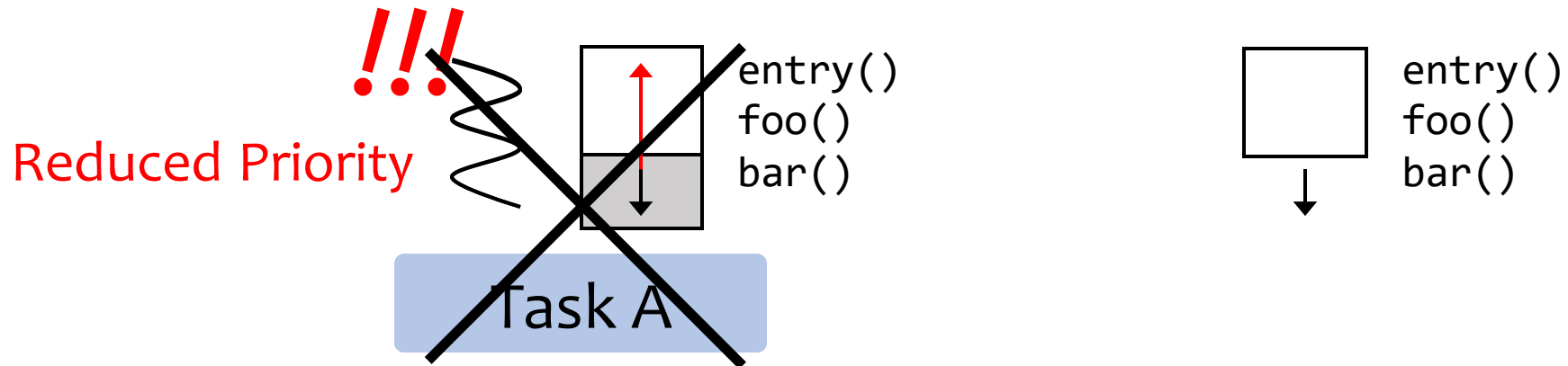
# RCB Supports Recovery from Panic

- Reliable Computing Base (RCB)
- Can recover panics from components above RCB



# Acceleration by Concurrent Task Restart

- Start a cloned task from entry() while unwinding the panicked one.
- Unwinding uses otherwise idle CPU time.



# Restartable Task Abstraction

Application  
Developer  
Supply

```
pub fn create_restartable_task<F, A>(
    entry_closure: F, entry_argument: A,
    stack_size: usize, priority: u8) where
F: FnOnce(A) + Clone + Send + Sync + 'static,
A: Clone + Send + Sync + 'static, { ... }
```

**Clone**: safe to duplicate

**Send**: safe to move across task context

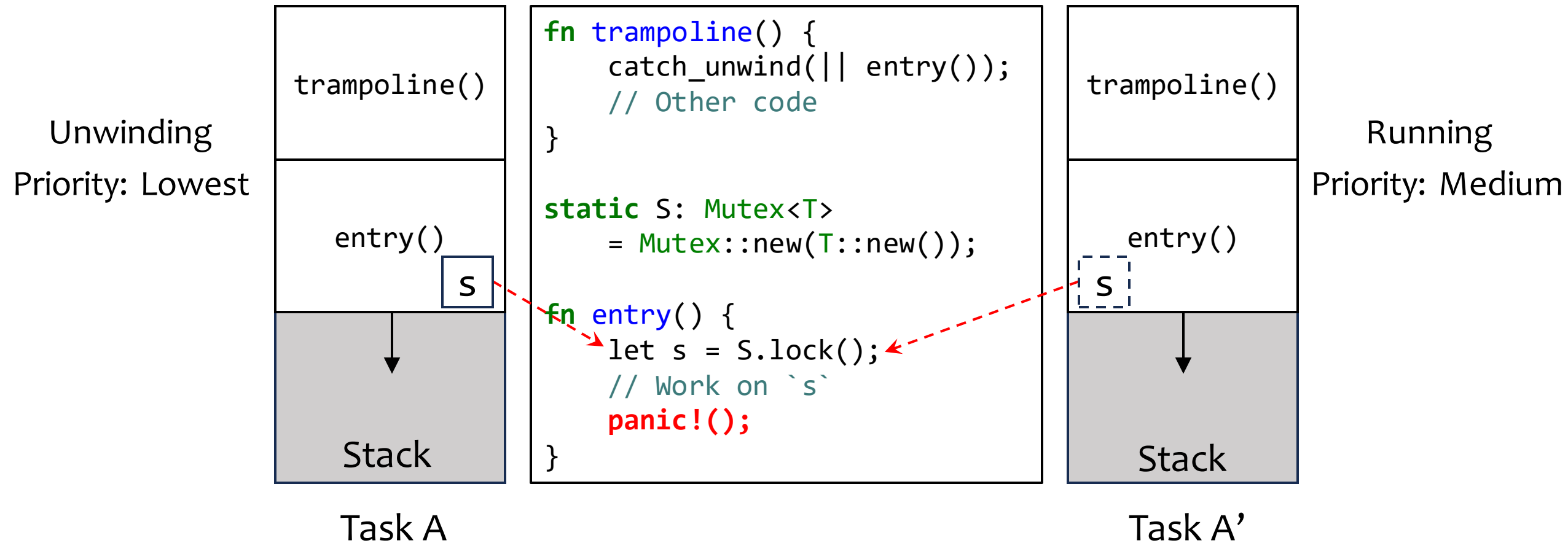
**Sync**: safe to access from multiple task context

# Restartable Task Abstraction – Continued

```
fn restartable_task_entry_trampoline<F, A>(
  entry_closure: &F, entry_argument: &A)
where ..., {
  catch_unwind_with_arg(
    entry_closure.clone(),
    entry_argument.clone());
}
```

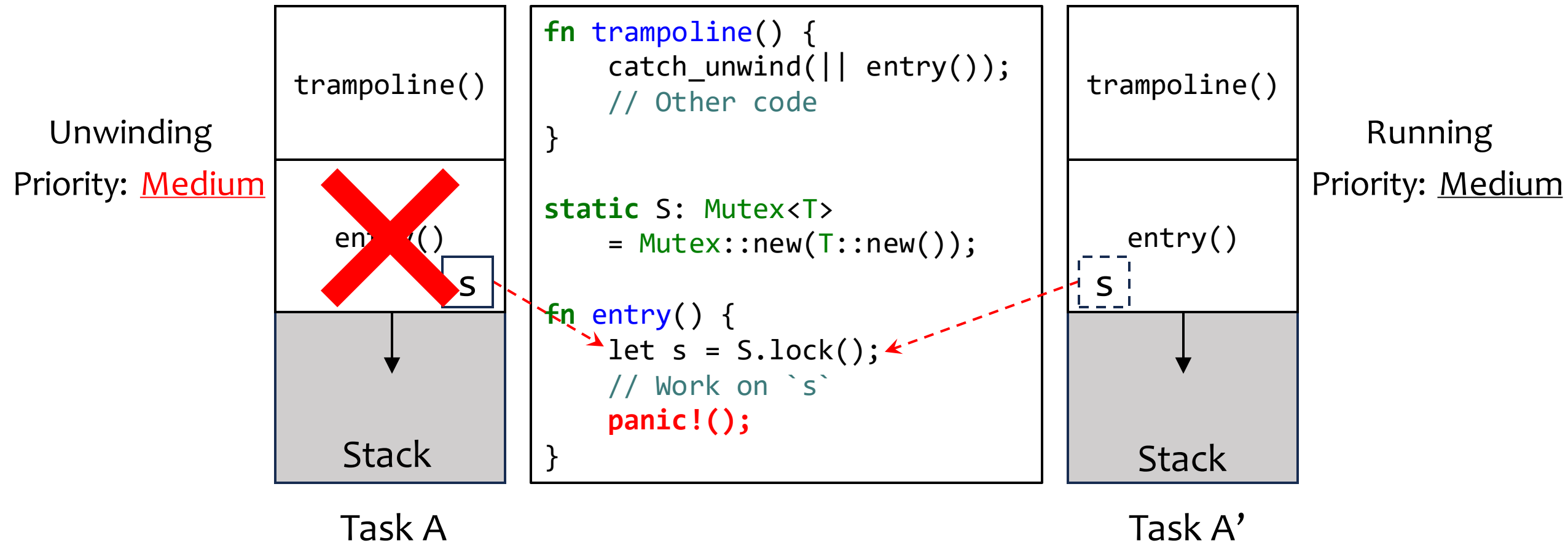
OS catches the panic outside of task's entry function.

# Priority Inheritance During Unwinding

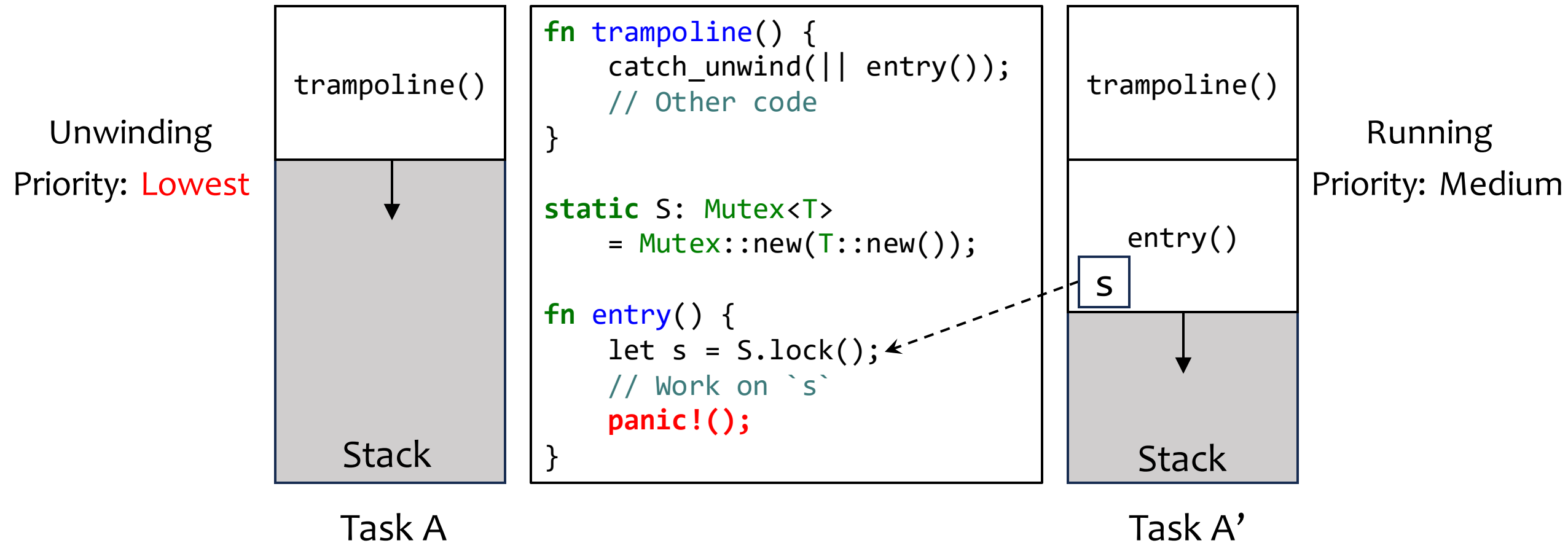




# Priority Inheritance During Unwinding



# Priority Inheritance During Unwinding



# Rust Facilitates Concurrent Task Restart

- Rust precludes race conditions.
  - Between panicked and restarted tasks, accessing static variables.
  - Safe Rust disallows mutable static variables and requires `Sync` trait.
  - → Must use `Mutex` around mutable static variables or atomic types
  - → `Mutex` priority inheritance works as normal

# Rust Facilitates Concurrent Task Restart

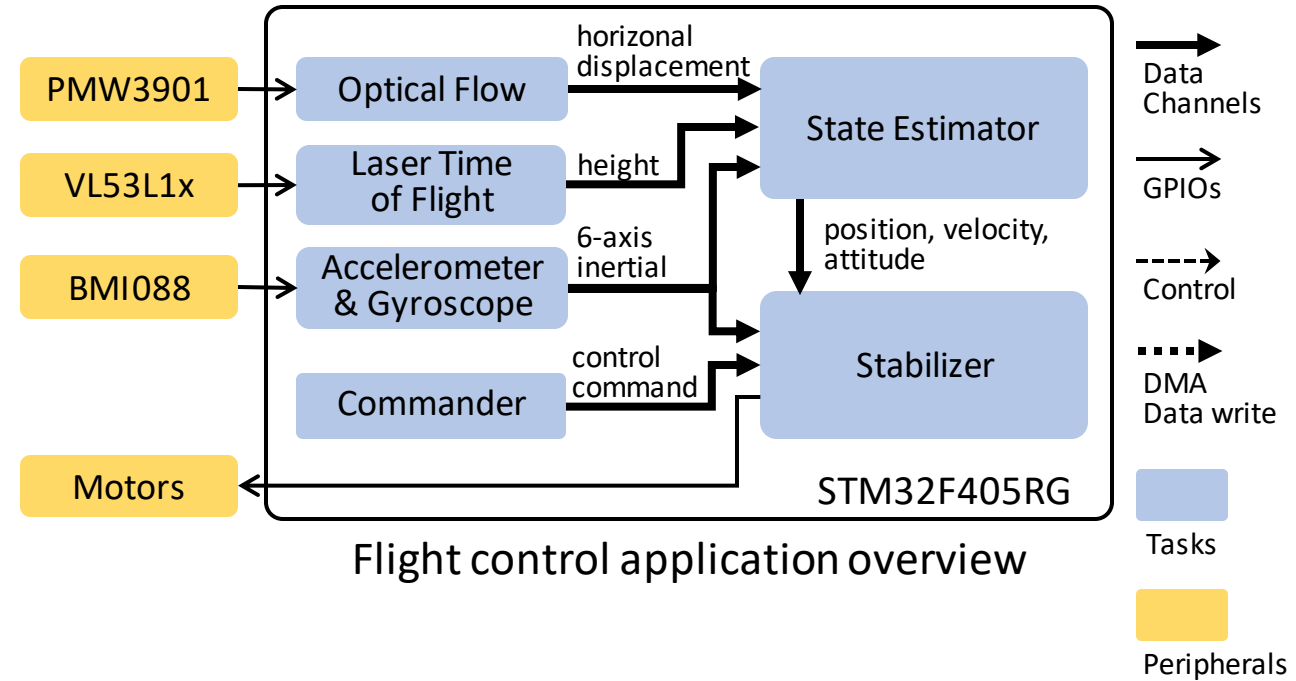
- Rust precludes race conditions.
  - Between panicked and restarted tasks, accessing static variables.
  - Safe Rust disallows mutable static variables and requires Sync trait.
  - → Must use Mutex around mutable static variables or atomic types
  - → Mutex priority inheritance works as normal
- Rust disambiguates fatal exception.
  - Rust panic is always fatal.
  - C++ exception not always fatal: `std::stoi()` throws `std::invalid_argument()`.
  - Concurrent restart only makes sense for fatal exception.

# Evaluation with a Flying Drone

An Example of Soft Real-time System

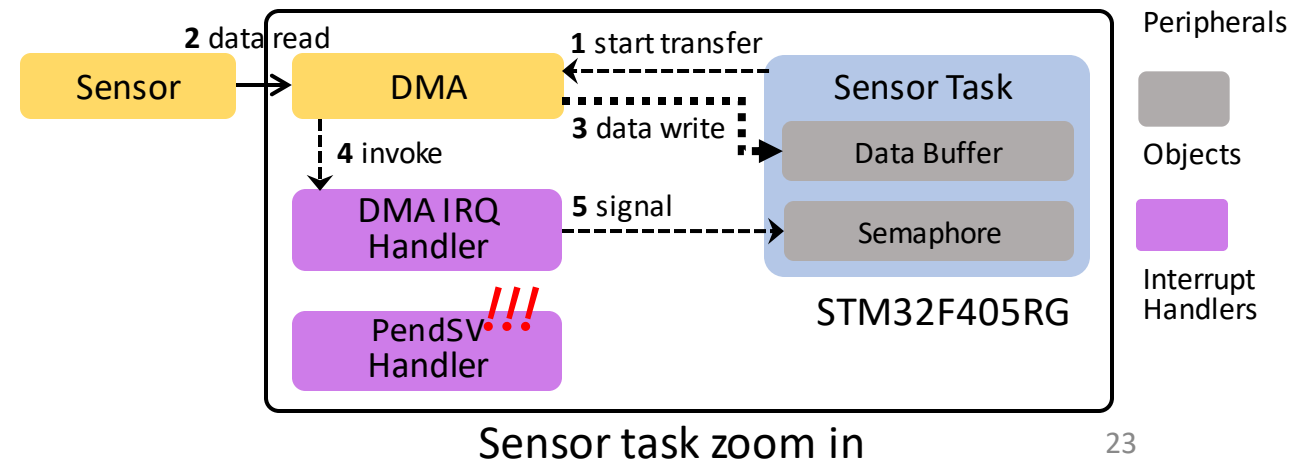
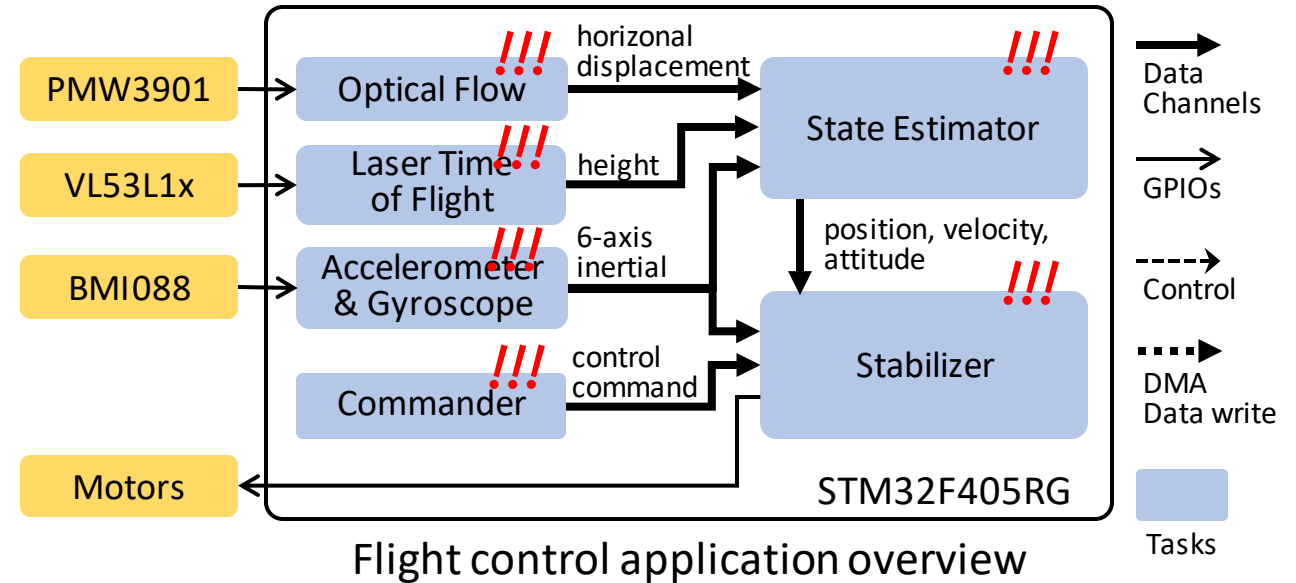
# Running Flight Control Application on Hopter

- Crazyflie 2.1
- COTS miniature drone
- Originally with FreeRTOS



# Can Sustain Panics from Task & IRQ Handler

- Put `panic!()` in task's code
- Put `panic!()` in PendSV handler



## Without Panic



## With Panic





# Price for unwinding

- 2.6% more CPU usage after enabling unwinding
  - Precluded compiler optimizations
- 26.0% storage overhead
  - Unwinder logic, landing pads, exception table

# Conclusions

- Panic recovery via unwinding is feasible on embedded systems.
- Soft real-time constraint can be met.
- Rust reduces unwinder complexity.
- Rust facilitates concurrent task restarts.
- 2.6% CPU overhead, 26.0% storage overhead when flying a drone.